

BitBlaze: Binary Analysis for Computer Security

Dawn Song

***Computer Science Dept.
UC Berkeley***

Malicious Code---Critical Threat on the Internet

- **Diverse forms**
 - Worms, botnets, spyware, viruses, trojan horses, etc.
- **High prevalence**
 - CodeRed Infected 500,000 servers
 - 61% U.S. computers infected with spyware [National Cyber Security Alliance06]
 - Millions of computers in botnets
- **Fast propagation**
 - Slammer scanned 90% Internet within 10 mins
- **Huge damage**
 - \$10billion annual financial loss [ComputerEconomics05]

Defense is Challenging

- **Software inevitably has bugs/security vulnerabilities**
 - Intrinsic complexity
 - Time-to-market pressure
 - Legacy code
 - Long time to produce/deploy patches
- **Attackers have real financial incentives to exploit them**
 - Thriving underground market
- **Large scale zombie platform for malicious activities**
- **Attacks increase in sophistication**
- **We need more effective techniques and tools for defense**
 - Previous approaches largely symptom & heuristics based

The BitBlaze Approach

- **Semantics based, focus on root cause:**

Automatically extracting security-related properties from binary code (vulnerable programs & malicious code) for effective defense

- **Automatically create high-quality detection & defense mechanisms**
 - Automatic generation of vulnerability signatures to filter out exploits
 - Automatic detection and classification of malware
 - » Spyware, keylogger, rootkit, etc.
 - » Automatic detection of botnet traffic
- **Able to handle binary-only setting**

Binary Analysis: Imperative & Challenging

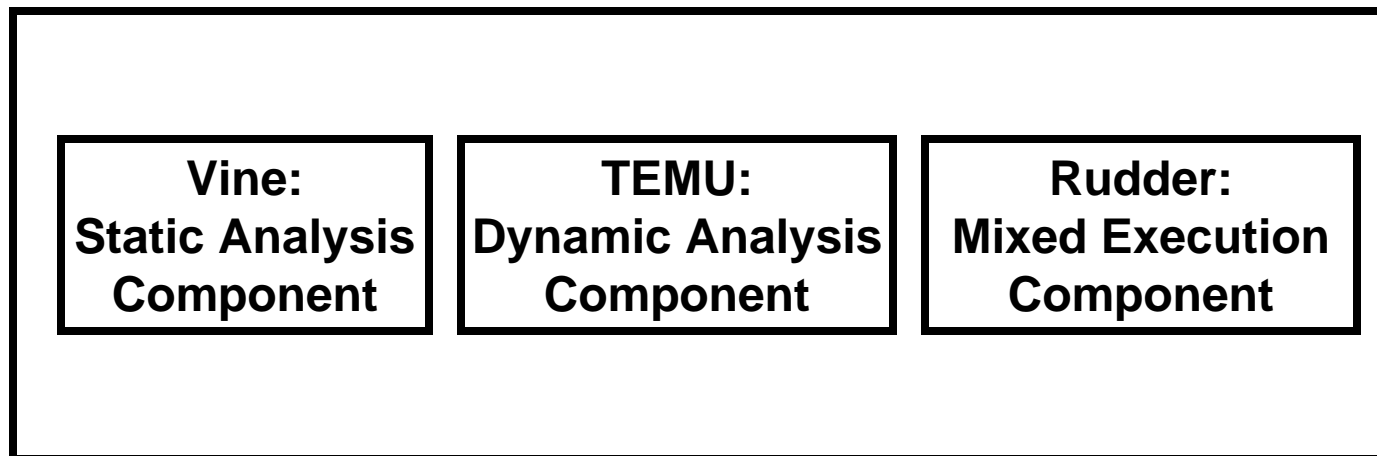
- **Binary analysis is imperative**
 - Source code is often unavailable
 - » COTS programs
 - » Malicious code
 - Binary is truthful
- **Binary analysis is challenging**
 - Lack higher-level semantics
 - » Even disassembling is non-trivial
 - Malicious code may obfuscate
 - » Code packing
 - » Code encryption
 - » Code obfuscation & dynamically generated code
- **Need techniques & tools to address these issues**

The BitBlaze Vision & Research Foci

- 1. Design and develop a unified binary analysis platform for security applications**
 - Identify & cater common needs of different security applications
 - Leverage recent advances in program analysis, formal methods, binary instrumentation/analysis techniques to enable new capabilities
- 2. Introduce binary-centric approach as a powerful arsenal to solve real-world security problems**
 - COTS vulnerability analysis & defense
 - Malicious code analysis & defense
 - Other security applications

The BitBlaze Binary Analysis Platform

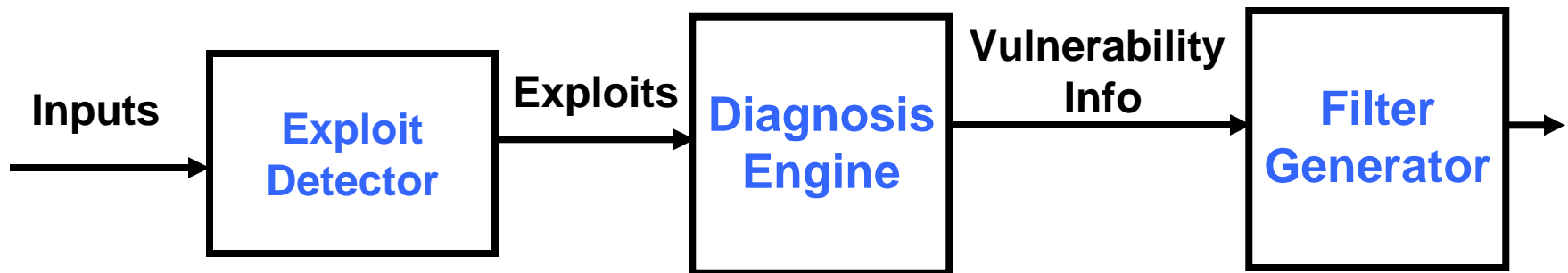
- **A unique infrastructure:**
 - **Novel fusion of static, dynamic analysis techniques, and formal analysis techniques such as symbolic execution**
 - **Vine: accurate static analysis using Vinell (Intermediate Language)**
 - **TEMU: whole-system, fine-grained, symbolic emulation system**
 - **Rudder: automatic exploration of program execution space**



BitBlaze Binary Analysis Platform

BitBlaze in Action: Addressing Security Problems

- **Effective new approaches for diverse security problems**
 - Over dozen projects
 - Over 12 publications in security conferences
- **Exploit detection, diagnosis, defense**



- **In-depth malware analysis**
- **Others:**
 - Reverse engineering
 - Deviation detection [Best Paper Award]
 - Semantic binary diff

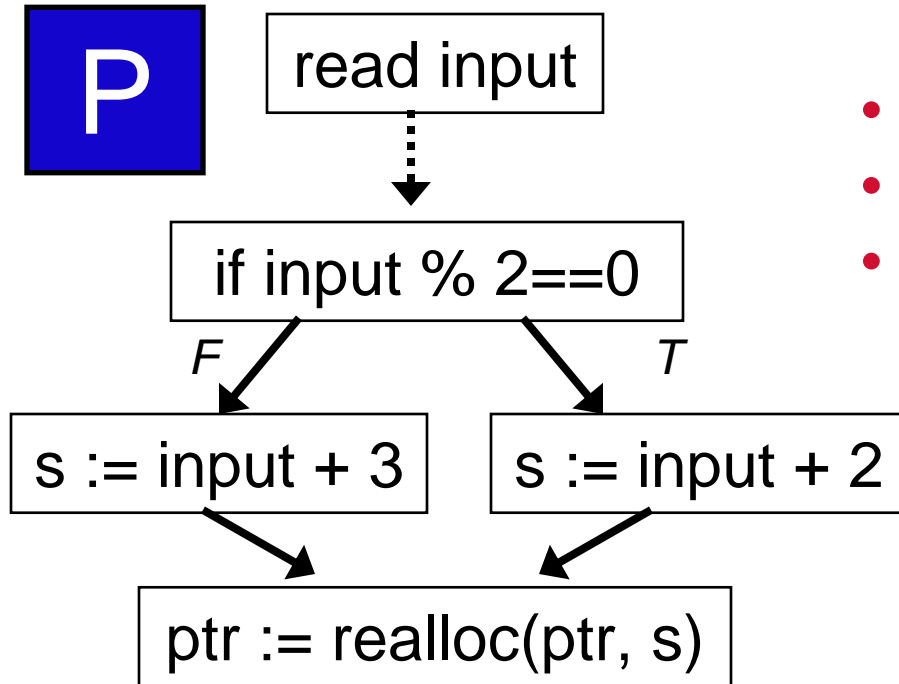
Talk Outline

- **Motivating security applications**
 - Automatic patch-based exploit generation
- **Components**
 - Vine: VineIR, static analysis on VineIR
 - TEMU: whole-system, fine-grained, symbolic emulation system
 - Rudder: automatic execution space exploration
- **Future directions and conclusion**

Automatic Patch-based Exploit Generation

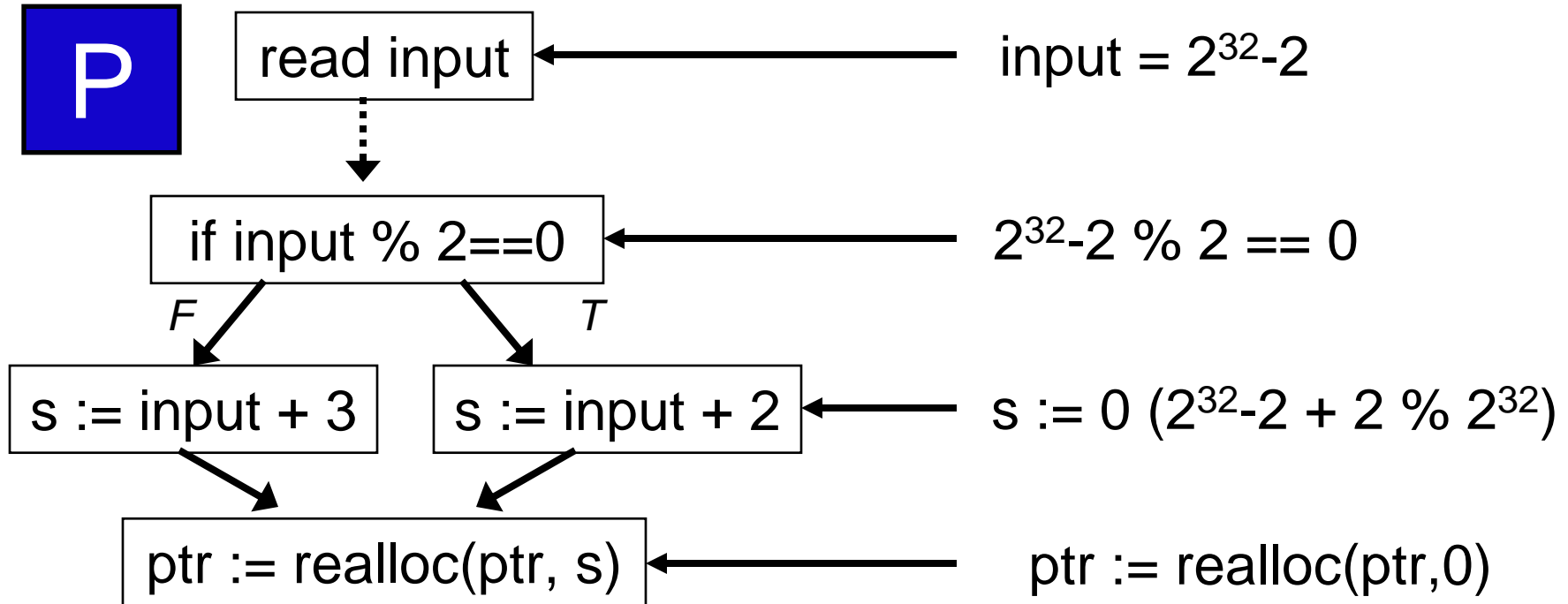
- **Given vulnerable program P, patched program P', automatically generate exploits for P**
- **Why care?**
 - **Exploits worth money**
 - » **Typically \$10,000 - \$100,000**
 - **Know thy enemy**
 - » **Security of patch distribution schemes?**
 - **Patch testing**

Running Example



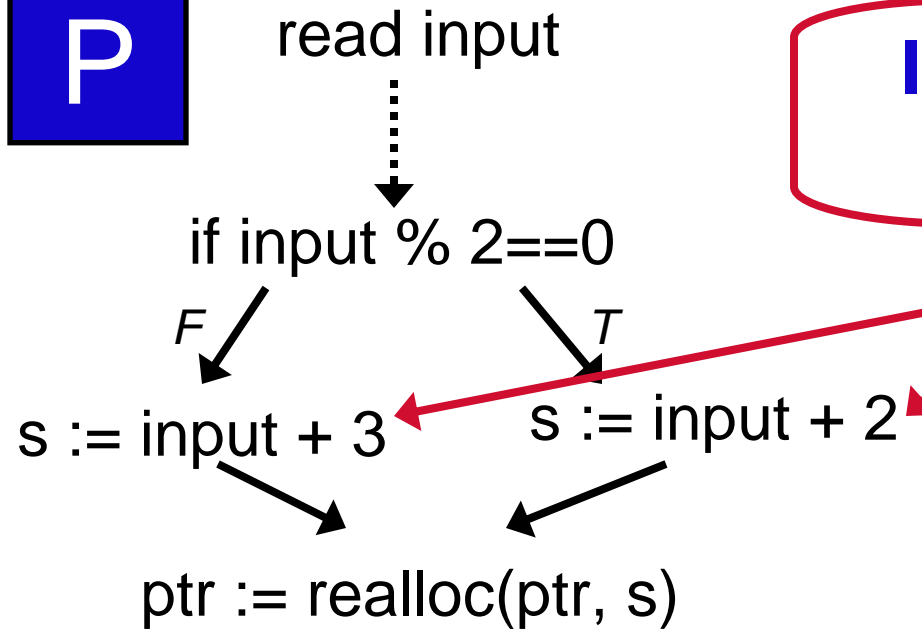
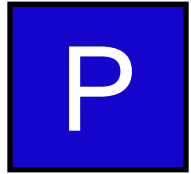
- All integers unsigned 32-bits
- All arithmetic mod 2^{32}
- Motivated by real-world vulnerability

Running Example



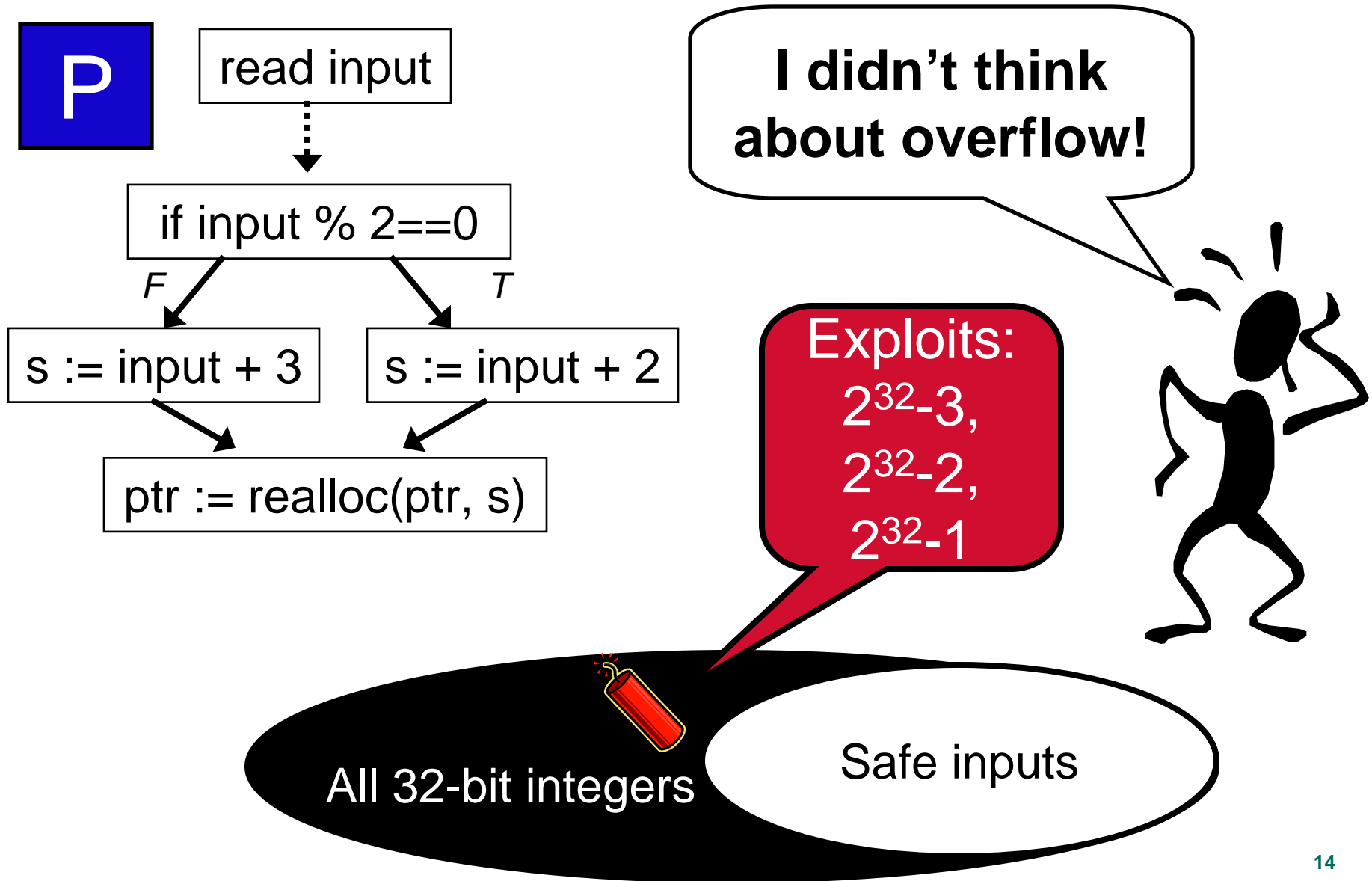
Using `ptr` is a problem

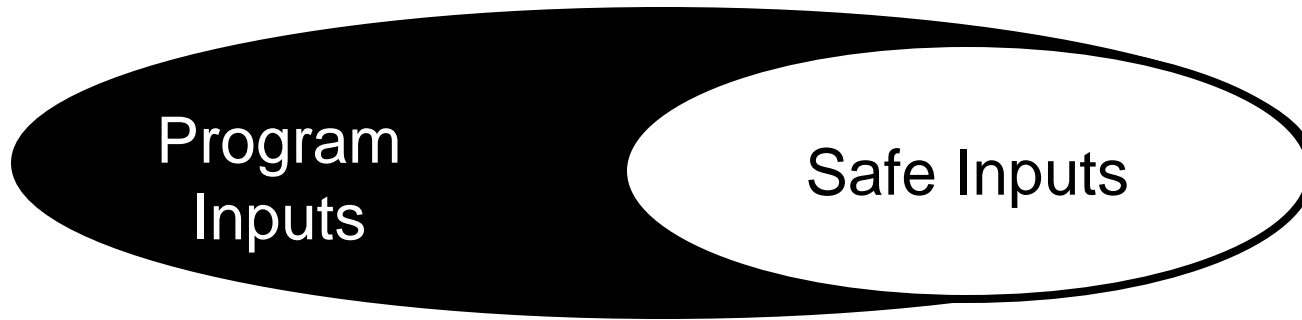
Running Example



**Integer Overflow when:
 $s < \text{input}$**

Running Example

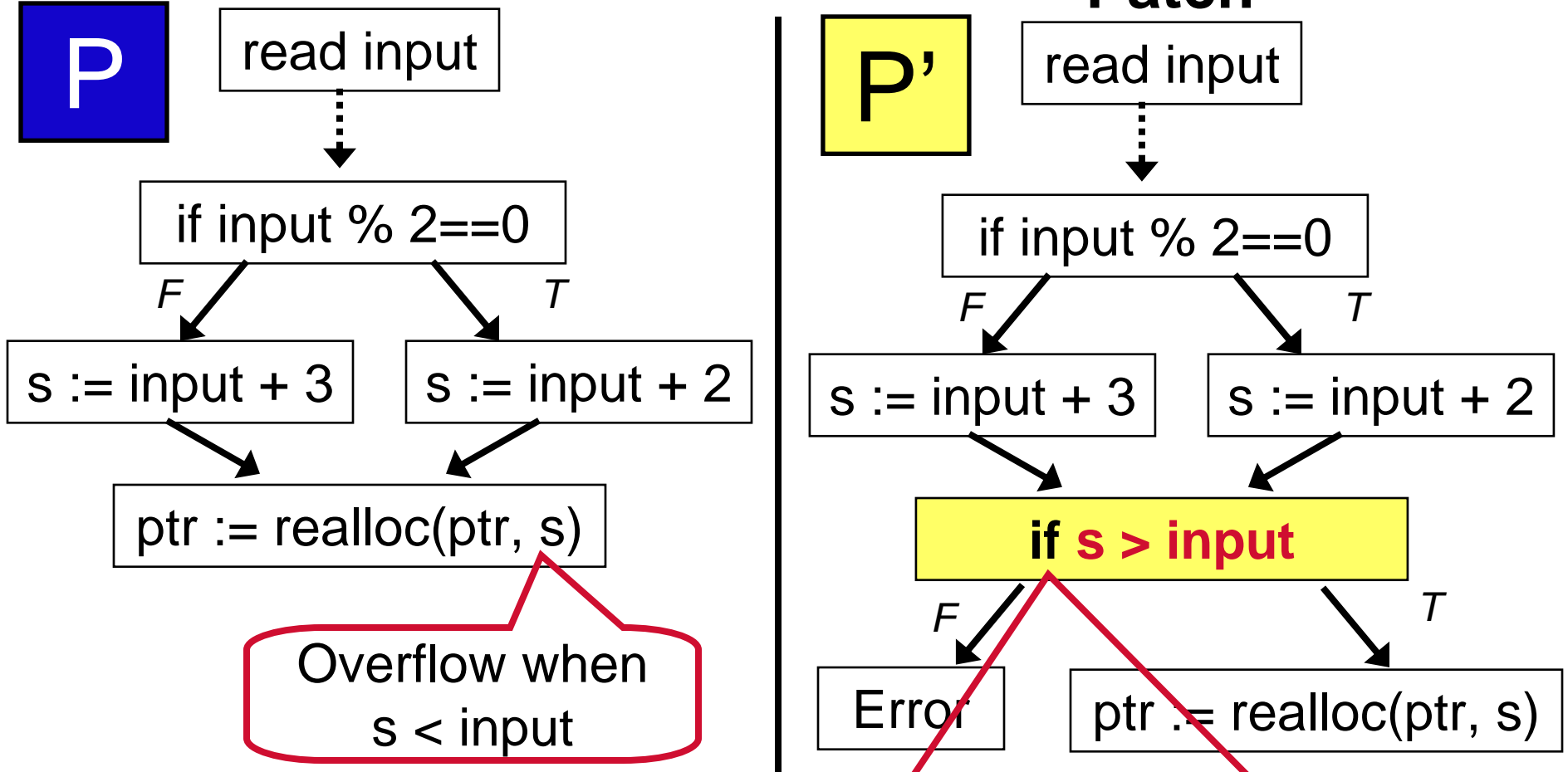




Input Validation Vulnerability

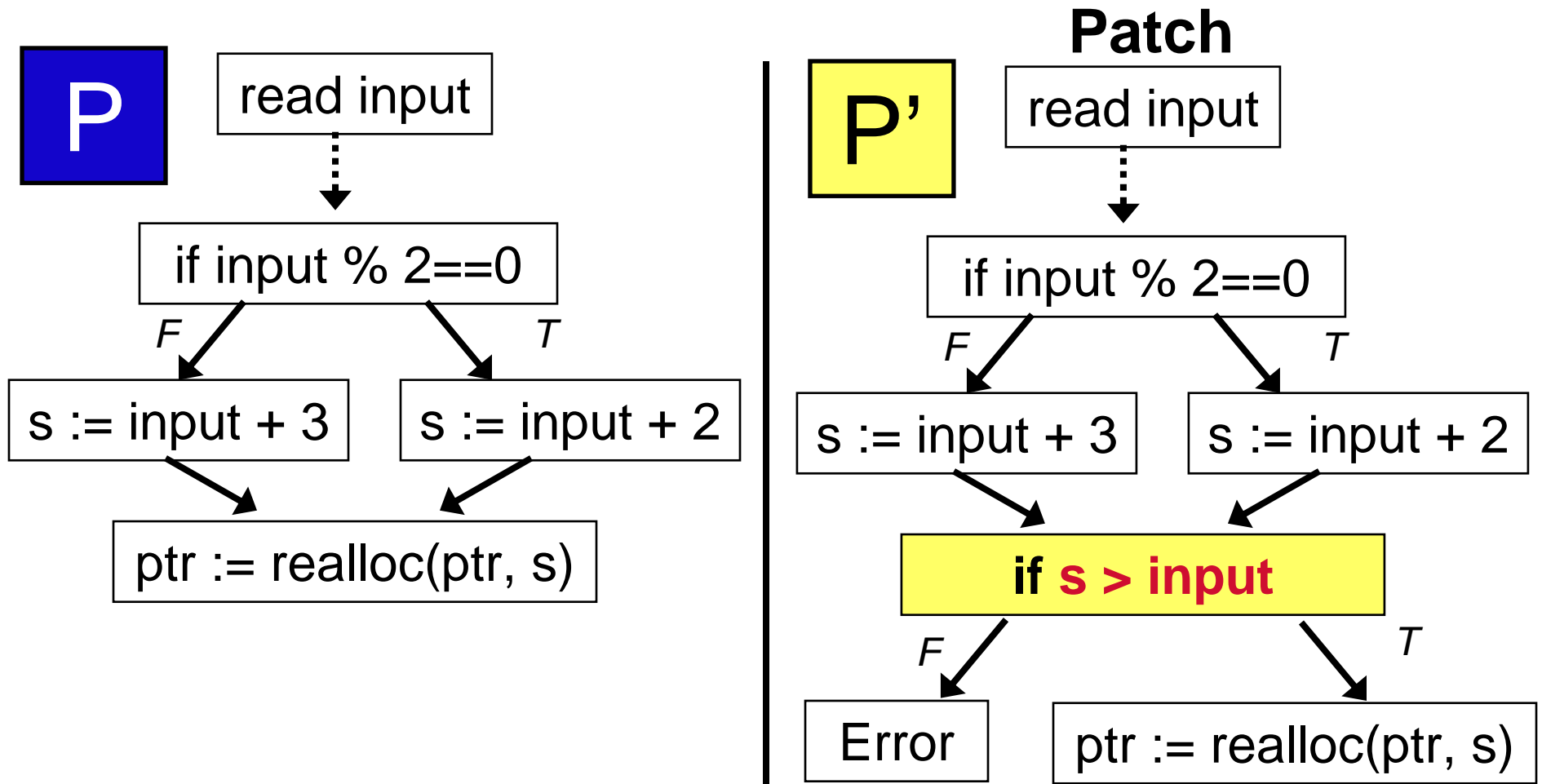
- **Programmer fails to sanitize inputs**
- **Large class of security-critical vulnerabilities**
 - “Buffer overflow”, “integer overflow”, “format string vulns”, etc.
- **Responsible for many, many compromised computers**

Patch



Patch leaks

1. **Vulnerability point** (where in code)
2. **Vulnerability condition** (under what conditions)

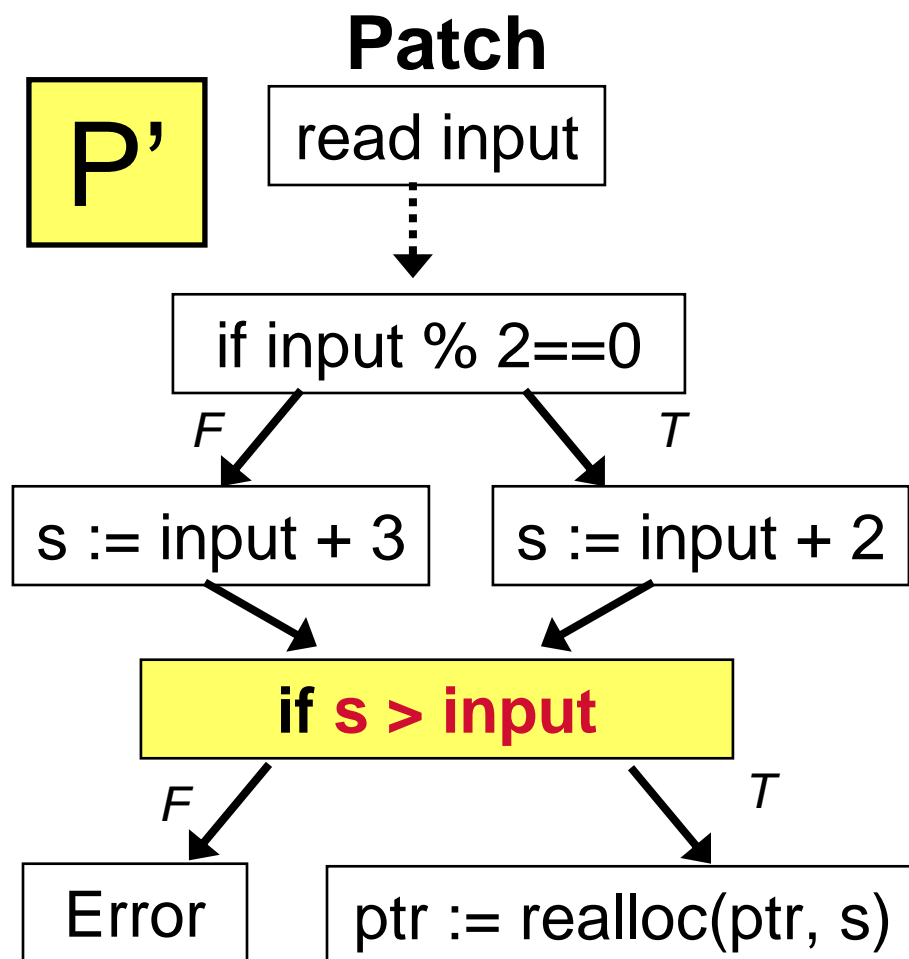


Exploits for P are inputs that fail vulnerability condition at vulnerability point
(s > input) = false

Our Approach for Patch-based Exploit Generation (I)

Exploit Generation

1. Diff P and P' to identify candidate vuln point and condition
2. Create input that satisfy candidate vuln condition in P'
 - i.e., candidate exploits
3. Check candidate exploits on P



Our Approach for Patch-based Exploit Generation (II)

- **Diff P and P' to identify candidate vuln point and condition**
 - Currently only consider inserted sanity checks
 - Use binary diffing tools to identify inserted checks
 - » Existing off-the-shelf syntactic diffing tools
 - » BinHunt: our semantic diffing tool
- **Create candidate exploits**
 - i.e., input that satisfy candidate vuln condition in P'
- **Validate candidate exploits on P**
 - E.g., dynamic taint analysis (TaintCheck)

Create Candidate Exploits

- **Given candidate vulnerability point & condition**
- **Compute Weakest Precondition over program paths**
 - Using vulnerability condition as post condition
 - Construct formulas representing conditions on input
 - » Whose execution path included
 - » Satisfying the vulnerability condition at vulnerability point
- **Solve formula using solvers**
 - E.g., decision procedures
 - Satisfying answers are candidate exploits

Different Approaches for Creating Formulas

- **Statically computing formula**
 - Covering many paths (without explicitly enumerating them)
 - Sometimes hard to solve formula
- **Dynamically computing formula**
 - Formula easier to solve
 - Covering only one path
- **Combined dynamic and static approach**
 - Covering multiple paths
 - Tune for formula complexity
- **Experimental results**
 - Different approach effective for different scenarios
- **Other techniques to make formulas smaller and easier to solve**

Experimental Results

- **5 Microsoft patches**
 - Mostly 2007
 - Integer overflow, buffer overflow, information disclosure, DoS
- **Automatically generated exploits for all 5 patches**
 - In seconds to minutes
 - 3 out of 5 have no publicly available exploits
 - Automatically generated exploit variants for the other 2
- **Diffing time**
 - A few minutes

Exploit Generation Results

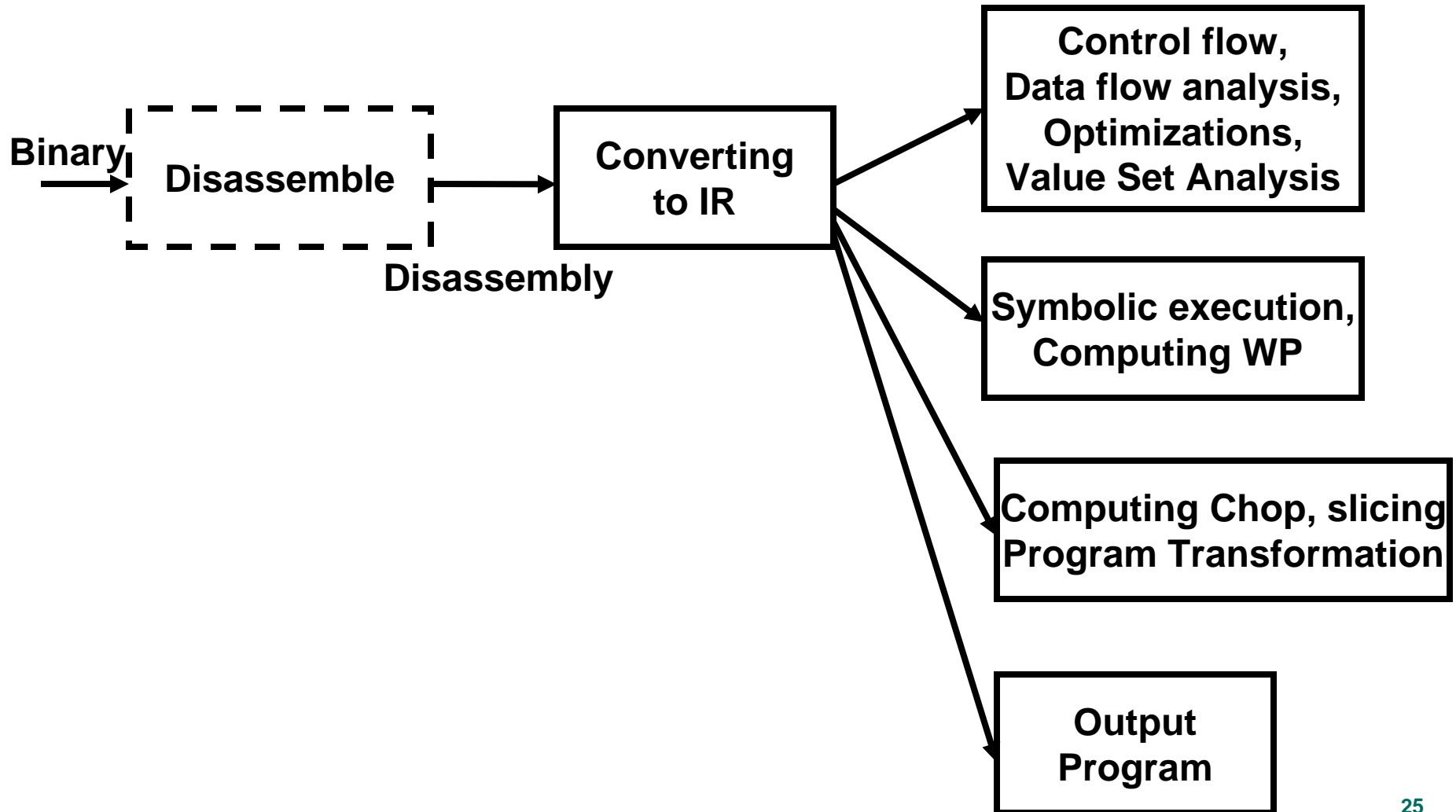
Time (s)	DSA_SetItem	ASPNet_Filter	GDI	IGMP	PNG
Dynamic Total	5.68	11.57	10.34	N/A	N/A
Formula	5.51	4.64	10.33	N/A	N/A
Solver	0.17	6.93	0.01	N/A	N/A
Static Total	83.47	N/A	26.41	N/A	N/A
Formula	2.32	N/A	4.99	N/A	N/A
Solver	81.15	N/A	21.42	N/A	N/A
Combined	11.51	N/A	29.07	13.57	104.28
Forumla	6.72	N/A	25.29	13.31	104.14
Solver	4.79	N/A	3.78	0.26	0.14

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Vine

- **Static analysis component**



Vine IR

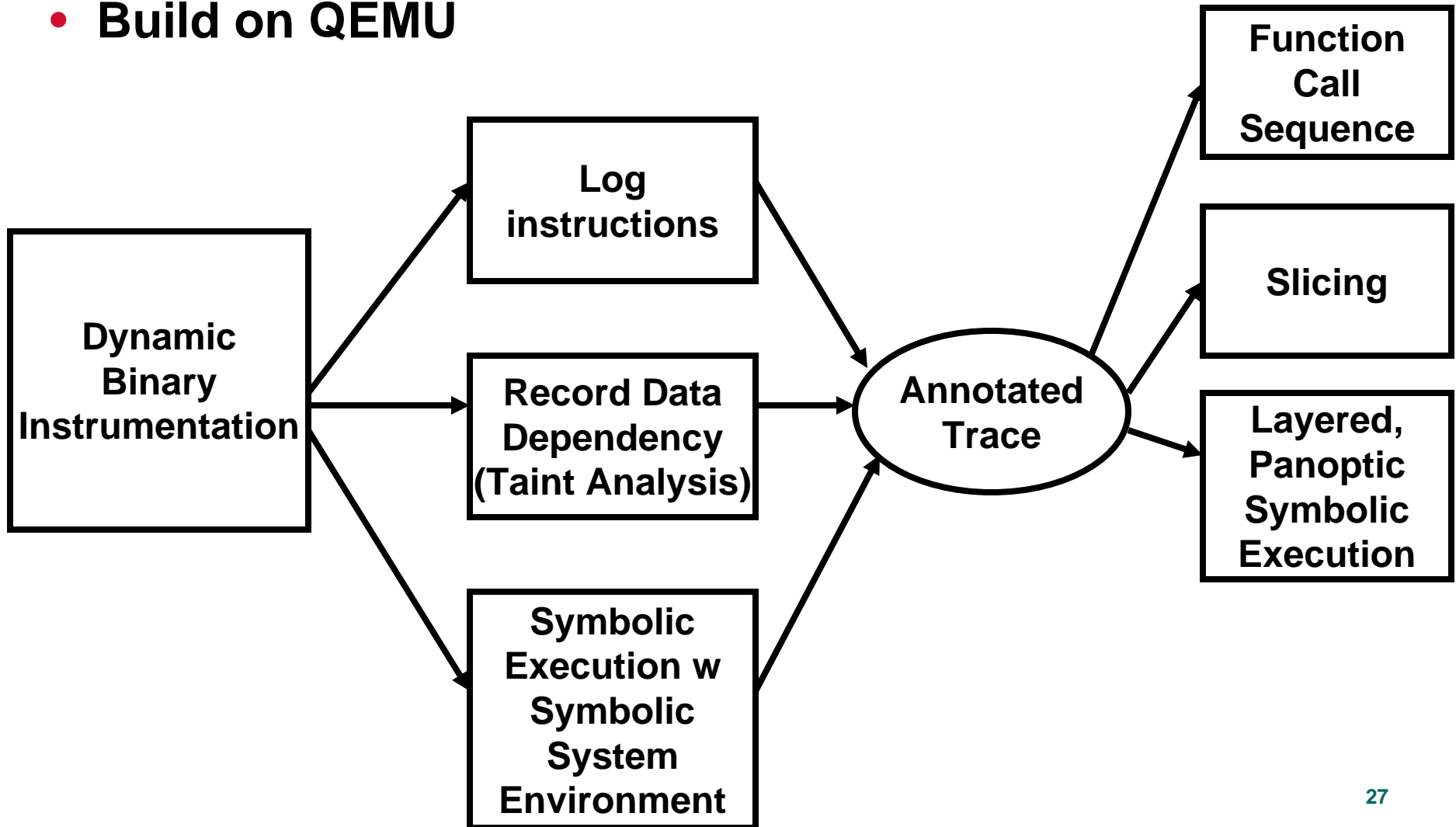
- **Simple RISC-like language, well-typed**

```
lval := exp
| goto exp
| if exp then goto exp1 else exp2
| return exp
| call exp
| assert exp
| special exp
| unknown (effects)
```

- **Handle x86, and ARM in progress**

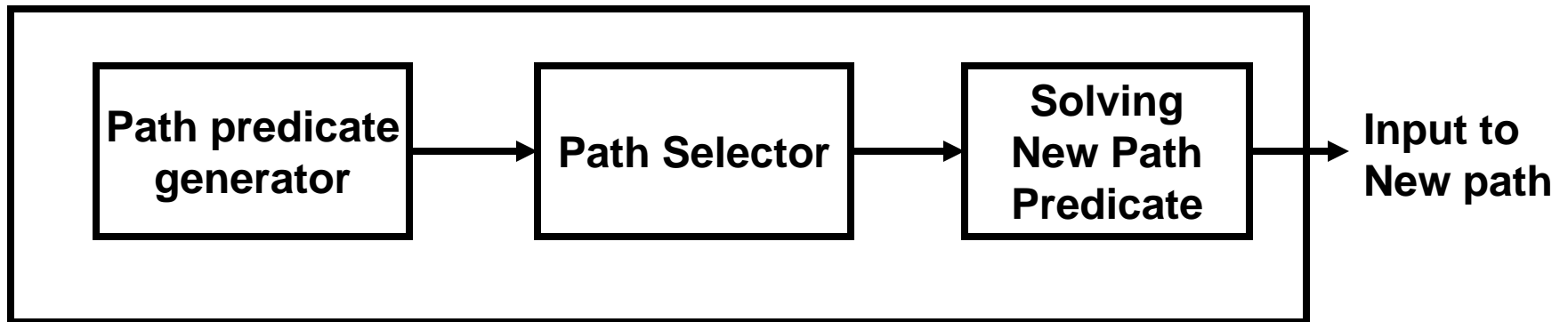
TEMU

- Work for both Windows & Linux, applications & kernel
- Build on QEMU



Rudder

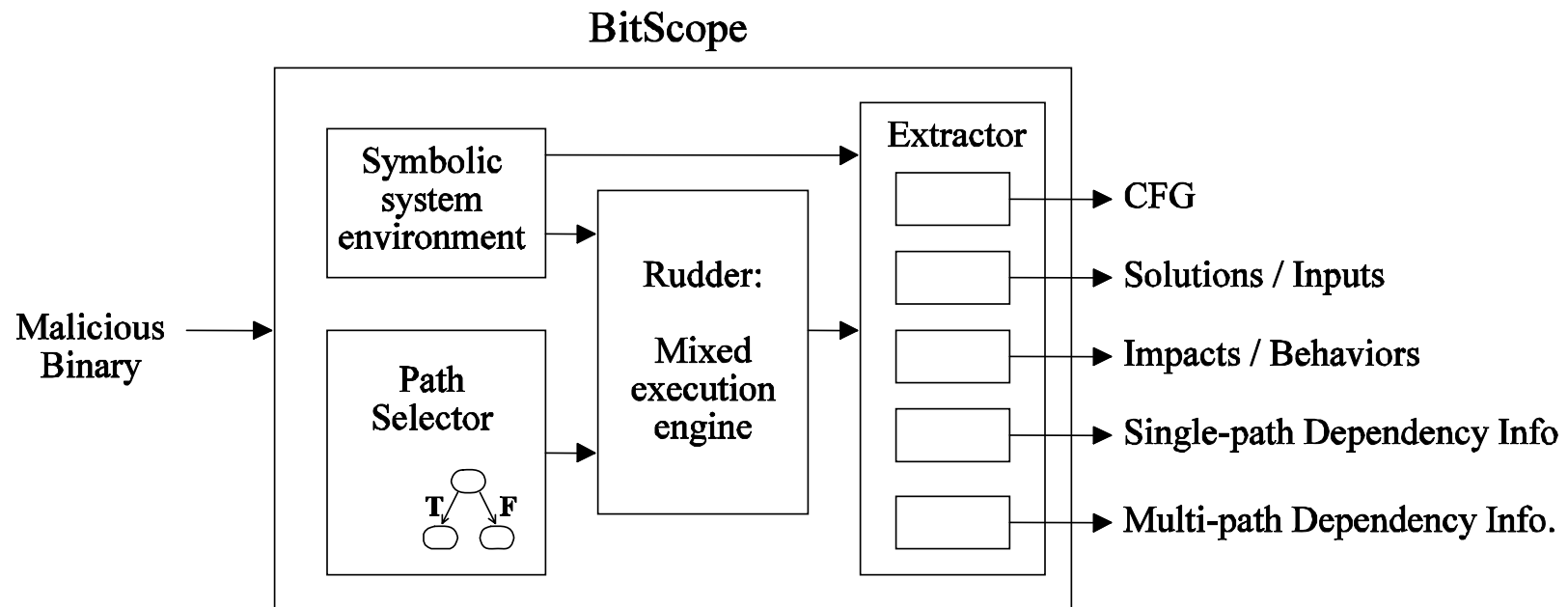
- **Compute path predicate**
- **Obtain new path predicate by reverting branches**
- **Solve path predicate to obtain new input to go down a different path**



Rudder

BitScope

- Built on top of TEMU & Rudder
- Work for packed code, self-encrypted code



BitScope: THE In-depth Malware Analysis infrastructure

- **Identify/analyze malicious behavior based on root cause**
 - Privacy-breaching malware: spyware, keylogger, backdoor, etc.
 - Malware perturbing system by hooking: rootkit, etc.
- **Understand how malware get into the system**
 - What mechanisms/vulnerabilities does it exploit
- **Explore hidden behavior, detect trigger-based behavior**
 - Automatically identifying botnet program commands, time bombs, etc.
- **Semantic & correlation analysis of malware input/output behavior**
 - Understanding the semantics of botnet program commands, etc.

Challenges

- **Performance & scalability for large programs**
- **Sample components we can take advantage of**
 - **Better identification of functions & resolution of indirect jumps**
 - » **Some of our VSA techniques may help**
 - **Better stack-walker**
 - **Binary aliasing analysis**
 - **More efficient binary instrumentation**

Conclusion

- **BitBlaze binary analysis platform**
 - A unique fusion of dynamic, static analysis & formal analysis (symbolic execution, WP, etc.)
- **Security Applications**
 - Vulnerability discovery, diagnosis, defense
 - In-depth malware analysis
 - Reverse engineering
 - Binary diffs
- **Components may support other applications**

Contact

- <http://bitblaze.cs.berkeley.edu>
- **dawnsong@cs.berkeley.edu**
- **BitBlaze team:**
David Brumley, Juan Caballero, Ivan Jager, Cody Hartwig,
Min Gyung Kang, Zhenkai Liang, James Newsome,
Pongsin Poosankam, Prateek Saxena, Heng Yin